

B2 Security Evaluation

(On-line version with links <http://www.multicians.org/b2.html>)

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Paragraphs by THVV unless noted.

This page still needs more information. Areas that are incomplete are marked **XXX**. Additions are welcome: please send mail to the editor.

Last update: 2015-01-12 14:19

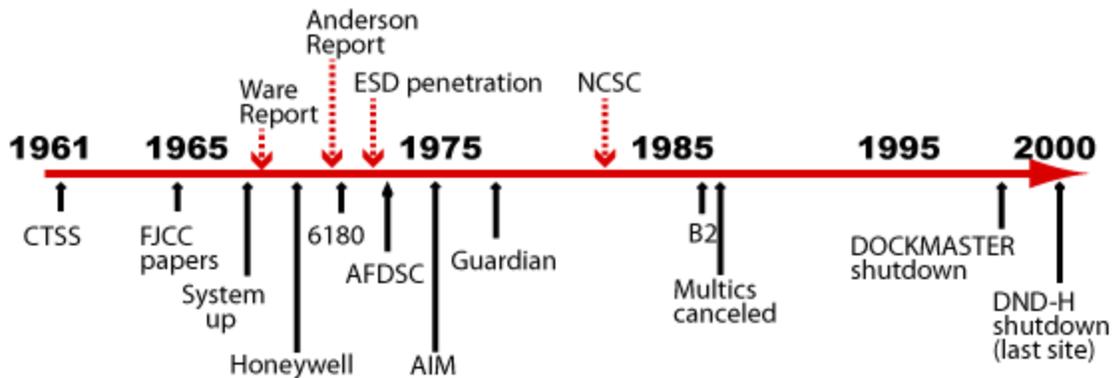
Introduction

When Multics was released in the early 1970s, it had a large collection of features designed to provide privacy and security; each organization that installed a Multics system chose how to use them to achieve its goals.

In the 1970s, the US military realized that their requirements for handling classified documents were not met by any computer system. Research led by US Air Force Major Roger Schell led to explicit models of US government document security requirements that military computer systems would have to meet. US DoD-funded projects built several demonstration systems and added security features to Multics.

In the 1980s, the US National Computer Security Center established a system security evaluation methodology, and the Multics team made changes to Multics and provided documentation, leading to the award of a class B2 rating.

This article describes the Multics security evaluation that led to the B2 rating in the mid 1980s, starting with the needs and context, describing the actual process, and discussing the results.



1. Early Computer Security

1.1 CTSS

MIT's CTSS, the predecessor of Multics, was designed in 1961 to provide multiple users with an independent share of a large computer. A major motivation for its security features was to ensure that the operating system and system accounting functioned correctly. Some CTSS security features were:

- In user mode, the CPU trapped I/O and privileged instructions.
- User login required a password (after March 1964).
- Users' files were protected from access by other users unless voluntarily shared.
- The system provided hierarchical administration and resource allocation.
- User program execution was done in a virtual machine with its own private unshared memory image.
- The memory bounds register in the CPU trapped user program attempts to modify the supervisor or other users' programs.

The original CTSS disk file system only supported file sharing within a work group. In September 1965, CTSS enhanced users' ability to share files by allowing *linking* to other users' files (if permitted), and added a "private mode" for files that limited use to only the file's owner.

1.2 Other Systems

A session at the 1967 Spring Joint Computer Conference on security stated the requirements for operating system security as they were understood then. It included papers by Willis Ware, Bernard Peters, and Rein Turn and Harold Petersen, and a panel discussion on privacy that included Ted Glaser. Butler Lampson's 1970 paper

on "Protection" [Lampson70] introduced the idea of an abstract access control matrix. [Peter Denning](#) and Scott Graham's 1972 paper "Protection - Principles and Practice" [Graham72] elaborated the matrix concept.

Almost as rapidly as computer systems implemented security measures in the 60s and 70s, the systems' protections were penetrated. [IBM76] Operating systems that were designed without security and had it added as an afterthought were particularly vulnerable.

1.3 Multics Security

The Multics objective to become a "[computer utility](#)" led us to a design for a tamper-proof system that allowed users to control the sharing of their files. Although [MIT Project MAC](#) was funded by the [Advanced Research Projects Agency](#) (ARPA) of the US Department of Defense, ARPA funding did not influence the design of Multics toward features specific to military applications. In the early 60s, [J. C. R. Licklider's](#) approach at ARPA's Information Processing Techniques Office (IPTO) was to support general improvements to the state of the computing art. One of the 1965 Multics conference papers, [Structure of the Multics Supervisor](#), stated "Multics does not safeguard against sustained and intelligently conceived espionage, and it is not intended to."

[JHS] Deployment of CTSS brought the discovery that users had a huge appetite to share programs and data with one another. Because of this discovery, Multics was designed to maximize the opportunity for sharing, but with control. The list of security features in Multics is long partly because controlled sharing requires elaborate mechanics.

[DMW] Security was an explicit goal of the Multics design and was continually reviewed as the implementation progressed. System components were designed to use the least level of privilege necessary.

When Multics was made available to users in the early 1970s, it had a large set of security-related features.

- In user mode, the CPU trapped I/O and privileged instructions.
- User login required a password.
- Users' files were protected from access by other users unless voluntarily shared.
- The system provided hierarchical administration and resource allocation.
- User program execution was done in [processes](#) with separate [virtual address spaces](#).

- Memory mapping registers in the CPU supported switching between virtual address spaces.
- Each process's virtual address space consisted of multiple segments.
- A segment could be part of more than one address space.
- Each process's address space was defined by a descriptor segment that defined segment access rights.
- Segments had different access permissions depending on the referencing process's descriptor segment.
- Segments had individual maximum sizes; out of bounds references caused a fault rather than accessing a different segment.
- Segments' accessibility depended on the process's ring of execution.
- In the Multics 6180 CPU architecture, the ring of execution was supported and checked by hardware; current ring number was kept in a CPU register.
- The CPU checked segment access rights and ring brackets on every instruction and generated a fault on unauthorized access.
- The 6180 CPU architecture provided cross-ring pointer checking.
- The operating system's privileged components executed in the most privileged ring.
- Segments were in-memory mappings of files in the file system.
- The file system defined segment accessibility with Access Control Lists (ACLs) and Ring Brackets.
- The system supported inner ring segments with extended access control permissions, such as message segments and mailboxes.
- Standard compiler output was non-writeable code (pure procedure).
- Standard compiler generated code used a stack segment for data; it was non-executable, as was the segment for static data.
- The PL/I string implementation kept track of the allocated size of a string.
- The call stack grew from lower addresses to higher, making it less likely that a buffer overrun would damage a return address pointer.
- Each process had a private directory for scratch storage (>pdd).

Mike Schroeder and Jerry Saltzer described the 6180's ring mechanism in [Schroeder71]. Jerry Saltzer's 1974 paper [Saltzer74] describes the design principles underlying Multics security, and ends with a candid description of weaknesses in the Multics protection mechanisms.

Multics developers thought that the combination of our hardware and software architectures, along with careful coding, could provide an adequate solution to the security problem. As mentioned in Saltzer's 1974 paper, we considered Multics to be "securable" at correctly managed installations. As it turned out, we had a long way to go.

1.4 US Military Computer Security Requirements

The US government had established levels of security classification (SECRET, TOP SECRET, etc) for documents before the introduction of computers. Initially, the military computed with secret data by dedicating a whole computer to each secure application, or by completely clearing the computer between jobs. Some military computer users desired multiple levels of security on a single computer, e.g. the US Military Airlift Command. [Schell12]

In 1967, a US government task force on computer security was set up, chaired by Willis Ware of RAND: the task force produced an influential report in 1970, known as the Ware report, "Security Controls for Computer Systems", that described the military's computer security needs. [Ware70] (The report was initially classified.)

1.4.1 GE/Honeywell Military Sales

The World Wide Military Command and Control System (WWMCCS) sale by GE/Honeywell Federal (\$46M Oct 1971, 35 mainframes, special version of GCOS, not multi-level) led to a significant expansion of GE's mainframe computer fabrication plant in Phoenix.

The Multics sale by GE/Honeywell Federal Systems Operation (FSO) in 1970 to Rome Air Development Center (RADC) at Griffiss Air Force Base in upstate New York did not depend on any security additions to Multics. This center did research on various security-related topics, but everything on the machine was at a single clearance level.

Honeywell FSO then bid Multics as a solution to a procurement by the US Air Force Data Services Center (AFDSC) in the Pentagon. AFDSC wanted to be able to provide time-sharing and to use a single system for multiple levels of classified information. AFDSC was a GE-635 GECOS customer already.

1.4.2 US Air Force Data Services Center Requirements

US Air Force Major Roger Schell obtained his PhD from MIT for work on Multics dynamic reconfiguration, and was then assigned to Air Force Electronics System Division (ESD) at L. G. Hanscom Field in Massachusetts in 1972. Roger had previously been a procurement officer for the USAF, and was familiar with the way that the US military specified and purchased electronic equipment. [Schell12]

Schell was assigned the problem of meeting the requirements of AFDSC, which wanted a computer system that could handle multiple levels of classified information

with controlled sharing. ESD contracted with James P. Anderson to evaluate the security of the GECOS machines for this purpose, and carried out penetration exercises for DoD on the GE-635. Ted Glaser was also involved in these studies.

Schell convened a panel that produced the Anderson Report, "Computer Security Technology Planning Study." Participants included James P. Anderson, [Ted Glaser](#), Clark Weissman, Eldred Nelson, [Dan Edwards](#), Hilda Faust, R. Stockton Gaines, and John Goodenough. (The report described buffer overflows and how to prevent them.) [Anderson72] In this report, Roger Schell introduced the concept of a *reference monitor*, implemented by a *security kernel*. [Schell12]

ESD funded projects at MITRE and Case Western Reserve University, directed by Schell, that described ways of modeling security. At MITRE, [Steve Lipner](#) managed [David Bell](#) and [Leonard LaPadula](#). [Bell12]

1.4.3 System Penetration Exercises

After he graduated from MIT (having done a thesis on Multics), the late [Paul Karger](#) fulfilled his Air Force ROTC commitment by joining ESD. Schell, Karger, and others carried out penetration exercises for DoD that showed that several commercial operating system products were vulnerable, including the famous [penetration of 645 Multics](#), [Karger74] which demonstrated problems in Multics hardware and software.

These studies showed Multics developers that despite architecture and care, security errors still existed in the system. Finding and patching implementation errors one by one wasn't a process guaranteed to remove all bugs. To be sure that 6180 Multics was secure, we needed a convincing proof of the security properties, and comprehensive testing. A proof of security would have to show that the system design implemented the security policy model, and verify that the code implementation matched the design. Furthermore, the Multics security kernel code in ring 0 was commingled with so much non-security code that verifying its correct operation was infeasible, as Saltzer had pointed out in his 1974 paper.

1.5 Multics MLS Controls Added for US Air Force

Roger Schell launched a project that would extend Multics to add multi-level security features called the [Access Isolation Mechanism](#)(AIM), in order to meet AFDSC requirements. The project proposal claimed that the architecture of Multics hardware and software protection supported the addition of mandatory checking to enforce the military document security classification policy. Funding for the project came from Honeywell, who hoped to sell computers to AFDSC, and ARPA IPTO, who wanted to

show military adoption of their research. Honeywell people included Earl Boebert, Charlie Spitzer, Lee Scheffler, Jerry Stern, Jerry Whitmore, Paul Green, and Doug Hunt.

AIM associated security classification levels and categories with objects on the computer, and security authorization levels and categories with users. AIM enforced mandatory access control as opposed to the discretionary access control provided by ACLs, and implemented the Star Property security policy specified by Bell and LaPadula. User processes had to have an authorization that was greater than or equal to the classification of data in order to operate on it (simple security), and information derived from a higher classification could not be written into a lower classification object (star property). No changes were made to limit covert timing channels.

Some specific features of the Multics AIM implementation were:

- Restrictions on the file system hierarchy were enforced: (1) files in a directory must have the same level as the directory. (2) directories contained in a directory must have the same level as the directory, or higher.
- Disk quota could not be moved between security levels, to avoid a covert storage channel.
- System building and tool dependencies were organized to create a secure supply chain.
- Changes were made to the answering service, device labeling, logging, I/O daemon queueing and labeling.
- New Storage System adaptations allowed the specification of a mandatory access range for file system devices.
- MITRE built an AIM test suite.

AFDSC eventually installed six Multics mainframe systems, two of which stored multi-level information. The systems that processed classified information were used in a secure environment where all users had clearances.

1.6 US Military Computer Security Projects in the Mid 1970s

In the mid 1970s, several other security efforts were funded by the US Department of Defense, through the Office of the Secretary of Defense (OSD), ARPA, and USAF.

1.6.1 Project Guardian

Project Guardian was a Honeywell/MIT/USAF project initiated by Roger Schell in 1975 to produce a high assurance minimized security kernel for Multics. [Schell73,

Schiller75b, Gilson76, HIS76, Schiller76a, Schiller76b, Schiller77] This work addressed problems described in Karger's security evaluation paper. [Karger74] The goal was to "produce an auditable version of Multics." [Schroeder78]

The implementation of Multics in the 60s and 70s preceded the ideas of a minimal security kernel; we included several facilities in ring 0 for efficiency, even though they didn't need to be there for security, such as the dynamic linker, included in the hardcore ring even though ring 0 never took a linkage fault. In 1974, Ring Zero comprised about 44,000 lines of code, mostly PL/I, plus about 10,000 lines of code in trusted processes.

Several design documents were produced by a Honeywell Federal Systems team led by Nate Adleman on how to produce a Multics organized around a small security kernel. [Adleman75, Biba75, Adleman76, Stern76, Withington78, Woodward78, Ames78, Ames80, Adleman80] This entailed moving non-security functions out of ring 0. A team at MIT Laboratory for Computer Science, led by Prof. Jerry Saltzer, continued design work on this project. Participants included Mike Schroeder, Dave Clark, Bob Mabee, and Doug Wells. This activity led to several conference papers. [Schell74, Schroeder75, Schroeder77, Janson75, Janson81]

Several MIT Laboratory for Computer Science TRs described individual projects for moving non-security features out of Ring Zero: see the References. These improvements to Multics never became part of the product. The prototype implementations looked like they would shrink the size of the kernel by half, but would have a noticeable performance impact. To adopt these designs into Multics and solve the performance problems would have required substantial additional effort by Multics developers.

1.6.2 SCOMP

SCOMP was a Honeywell project initiated by Roger Schell at ESD to create a secure front-end communications processor for Multics. [Gilson76, Fraim83] Roger Schell described the initial architecture, which was strongly influenced by Multics. The target customer was the US Air Naval Electronics System Command. Schell obtained funding for the project from ARPA and from small allocations in various parts of DoD. The project ran from 1975-77. SCOMP included formally verified hardware and software. [Gligor85, Benzel84] The hardware was a Honeywell Level 6 minicomputer, augmented with a custom Security Protection Module (SPM), and a modified CPU whose virtual memory unit obtained descriptors from the SPM when translating virtual addresses to real. All authorization was checked in hardware. SCOMP had a security kernel and 4 hardware rings. The SCOMP OS, called STOP, was organized into layers: the lowest layer was a security kernel written in Pascal that

simulated 32 compartments. STOP used 23 processes to provide trusted functions running in the environment provided by the kernel. A Formal Top Level Specification (FTLS) for the system's functions, written in the SRI language SPECIAL, was verified to implement DoD security rules.

1.6.3 Other Secure Systems

Additional research related to military computer security was carried out at other institutions, many by people who had been associated with Multics:

- MITRE: PDP-11 security kernel, [Lee Schiller](#) [Schiller75a] (Roger Schell architect, managed by Steve Lipner)
- SRI: PSOS, [Peter Neumann](#) (BTL Multics architect), [Rich Feiertag](#) (MIT Multics developer)
- Ford Aerospace: KSOS, Peter Neumann, 1978, Boyer-Moore/Modula 2, ran on DEC PDP-11/45 and PDP-11/70
- Secure Ada Target, Earl Boebert (HIS Multics developer), PSOS based, evolved into LOCK
- UCLA: [Gerry Popek](#) (Harvard thesis based on Multics) secure UNIX 1979

1.7 US Military Computer Security Projects at the End of the 70s

Air Force ESD was directed to discontinue computer security research by the Office of the Secretary of Defense in 1976. This led to several consequences:

- The Multics Guardian project was canceled, without shrinking ring 0. Some final reports were published. [HIS77]
- Roger Schell was reassigned to the US Air War College and then to the Naval Postgraduate School.
- Karger and Lipner went to DEC, to work on Secure VMS.

Several projects did continue in the late 1970s.

- DoD Computer Security Initiative, 1978.
- NBS workshop 1978 (Ted Lee, Peter Neumann, Gerry Popek, Steve Walker, Pete Tasker, Clark Weissman) Miami
- Air Force summer study at Draper
- MITRE work, Nibaldi report, introduced TCB [Nibaldi79]

2. US National Computer Security Center

The US Department of Defense Computer Security Center (DoDCSC) was established by DoD Directive 5215.1, dated October 25, 1982, as "a separate and unique entity within the NSA." Roger Schell credits Steve Walker as an important influencer in DoD's creation of this agency. The Center was tasked with "the conduct of trusted computer system evaluation and technical research activities" and the creation of an Evaluated Products List. NCSC addressed the US military's concern with the cost of acquiring and maintaining special-purpose computer systems with adequate security for classified information. [DOD82]

The late Melville Klein of NSA was appointed Director of the Center, and Col. Roger Schell Deputy Director. DoDCSC was renamed the National Computer Security Center (NCSC) in 1985. [Schell12] Among the 35 people initially assigned to the center were Dan Edwards and Marv Schaefer of NSA.

2.1 Theory of Security Evaluation

The basic idea behind NCSC's approach to security was that it was provided by a tamper-proof, inescapable *reference monitor*, as described by Roger Schell in the Anderson report. Abstractly, the reference monitor checked every operation by a *subject* on an *object*.

Based on work done in workshops and at MITRE, described above, the NCSC established that there should be several levels of computer system assurance, with different systems assured to different levels. The levels would be tiered by the threats that systems were exposed to: serious attacks, use only in a controlled environment, or unskilled opponent. Different levels of assurance would have different levels of testing. Systems that were assured to a higher level would give additional consideration to environmental factors, such as secure installation and maintenance.

2.2 Orange book

The Orange Book was published by the NCSC on 15 aug 1983 and reissued as DoD 5200.28-STD in December 1985. Sheila Brand of NSA was the editor. [NCSC83] Architects of this document were Dan Edwards, Roger Schell, Grace Hammonds, Peter Tasker, Marv Schaefer, and Ted Lee. Although it was explicitly stated in DODD 5215.1 that presence on the EPL should have no impact on procurement, over time military RFPs began to require features from the Orange Book.

The Orange Book defined security levels A1 (highest), B3, B2, B1, C2, C1, and D. Roger Schell wrote an IEEE conference paper describing the ratings and the evaluation methodology. [Schell83]

Level B2, "Structured Protection," the level Multics aimed at, required formal documentation of the security policy; identifiable and analyzable security mechanisms; mandatory access control on all subjects and objects; analysis of covert storage channels and audit of their use; a trusted path to the system login facility; trusted facility management; and strict configuration management controls.

To attain B3, "Security Domains," Multics would have had to have the OS restructured to move non-security functions out of ring 0, as Project Guardian planned; provide additional auditing of security-relevant events; built-in intrusion detection, notification, and response; trusted system recovery procedures; significant system engineering directed toward minimizing complexity; a defined security administrator role; and analysis of covert timing channels.

A1, "Verified Design," was the highest rating. It required the same functions as B3, but mandated the use of formal design and verification techniques including a formal top-level specification, as done with SCOMP, and also formal management and distribution procedures.

3. Multics Orange Book Evaluation in the 1980s

Naturally Honeywell wanted Multics to be evaluated according to the Orange Book criteria, and hoped that a high rating would make Multics attractive to more customers. Many people in the security community felt that Multics had been the test bed for concepts that later went into the Orange Book, and expected that the evaluation would be straightforward.

As Marv Schaefer describes in "What was the Question" [Schaefer04], the Orange Book criteria turned out to be imprecisely stated, and required many sessions of interpretation. As Schaefer says, "the process that became part of NCSC-lore made *slow* look fast in comparison."

[Olin Sibert] Multics was widely considered to be the prototypical B2 system, and there was some effort made to ensure that the criteria (Orange book) were written so that it could pass--and that was clearly the intended outcome. There turned out to be some significant lapses, both in terms of product features (auditing of successful accesses; control of administrative roles) and development process (functional tests for security features; software configuration management; design documentation), but overall, the only thing that could have prevented a successful outcome was cancellation of the product itself. Fortunately, Honeywell waited until evaluation had just been completed to do that.

3.1 Steps in the Multics Evaluation

XXX Table of dates

The Multics evaluation was an early one, and I don't have documentation of the exact steps taken or their dates. When I get a copy of the FER, maybe this will describe what was done. (This process was later codified in NCSC-TG-002 [Bright Blue Book] Trusted Product Evaluation A Guide for Vendors.)

3.1.1 Honeywell Decision

XXX Presumably Honeywell had to decide to go for this rating. How did the project start, and when? Someone in FSO asked for it, who (or did NSA ask us?) Somehow this was deliberated on and agreed to, by whom? Somehow FSO asked Engineering for it and it was approved, by whom? A plan was made, by whom? Presumably they went through a budgeting exercise and found funds for it.

3.1.2 Pre-Evaluation Review

Before actual evaluation began, there were a few steps. In the Multics case, the NCSC may have solicited a letter from Honeywell, or Honeywell FSO Marketing may have initiated the process. Since Honeywell had already sold computers to the US DoD, some of the paperwork would already be in place.

The NCSC preferred real products that would actually be sold: they didn't want "government specials."

When the NCSC accepted the evaluation proposal, there would be a Memorandum of Understanding signed by Honeywell and NCSC.

XXX Were all these documents created? if so, are they saved anywhere? Who signed them on each side?

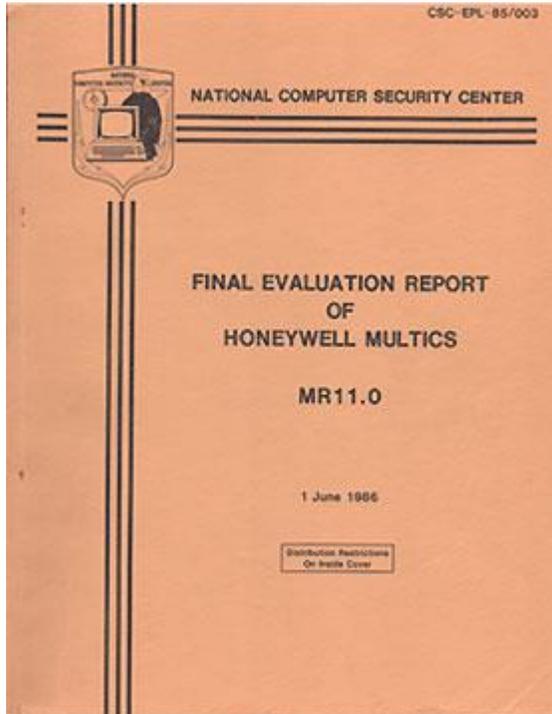
3.1.3 Initial Product Assessment

The usual next step was for a vendor to prepare a package of design documentation for its product and send it to the NCSC. In the early days of evaluations, this process was probably informal and iterative.

In the Multics case, the NCSC had tailored the Orange Book requirements for B2 with an eye to Multics.

[WOS] The Multics documentation "requirements" were the subject of lengthy negotiations between CISL and the evaluation team, and before Keith Loeper took ownership of the activity, I think I was probably responsible for most of that (working on the NCSC side).

3.1.4 Formal Evaluation



XXX who was the FSO project leader for B2 (worked for Earl Boebert)? who else in Marketing was involved?

the CISL effort for B2 must have been costly. what was the total cost, where did the funds come from, who approved it?

documentation was created as part of the B2 project. who figured out what docs were required?

there must have been some kind of schedule. did it take as long as it was expected to? did it cost as much?

was there opposition from Phoenix management, or other parts of Honeywell/Bull?

what other paperwork was agreed to between the NCSC and Honeywell?

did the NCSC release a Product Bulletin saying Multics was being evaluated for B2?

The goal of the evaluation phase was a detailed analysis of the hardware and software components of Multics, all system documentation, and a mapping of the security features and assurances to the Orange Book. This included actual use of Multics and penetration testing. These activities were carried out by an **Evaluation Team**, which

consisted of US government employees and contractors. The evaluators produced the evaluation report under the guidance of the senior NCSC staff.

The result of the evaluation phase was the release of the Multics Final Evaluation Report [NCSC85], on 1 June 1986, and an entry in the Evaluated Products List.

3.2 Honeywell

([THVV] I left the Multics group in 1981 and missed the whole B2 project, so I'll have to depend on others for the story.)

3.2.1 Honeywell Federal Systems

[WEB] The project was run out of the Arlington VA sales office on a day to day basis, dotted line reporting to me.

3.2.2 CISL

Honeywell's Cambridge Information Systems Laboratory (CISL) was the primary development organization for Multics.

People and roles:

- John Gintell -- manager of CISL.
- Michael Tague -- project mgr, wrote Bell and LaPadula MDD, Opus mgr
- Benson Margulies -- tech lead, wrote overview MDD
- Maria Pozzo -- wrote rcp MDD
- Olin Sibert -- wrote TCB MDD
- Keith Loepere -- tech lead, wrote initialization MDD, wrote SIGOPS journal paper on covert channels
- Gary Palter
- Ed Ranzenbach
- Ed Sharpe -- wrote several MDDs
- Eric Swenson -- wrote system and user control MDD
- Chris Jones -- wrote IOI MDD
- Mike Pandolfo -- wrote message seg MDD, rewrote ring-1 mail
- Robert Coren -- wrote traffic control MDD
- Melanie Weaver -- wrote runtime MDD
- Lynda Adams
- George Gilcrease -- wrote IO daemon MDD
- Ron Beatty
- John Gilson
- Allen Ball - student

- Wilson Wong

3.2.3 PMDC

Honeywell's Phoenix Multics Development Center (PMDC) was responsible for release testing and distribution, as well as development of some Multics subsystems.

People and roles:

- Frank Martinson
- Bonnie Braun
- Paul Dickson -- functional testing MDD
- Jim Lippard
- Bill May
- Rich Holmstedt
- Gary Dixon - MDD on DC, post B2 testing.
- Kevin Fleming

3.3 NCSC

3.3.1 NCSC Staff

[MS] I think you should recall from the powder-blue version of the TCSEC that was handed out without restriction at the NBS/DoDCSEC security conference in 1982 -- and I believe -- in print in the IEEE Oakland Symposium paper by Roger Schell introducing either the Center or the Criteria, that the Orange Book was written with specific worked examples in mind: RACF for what became C2 (but which turned out only to get a C1), B1 as the 'secure' Unix (labels with training wheels), B2 as AFDSC Multics, and A1 as KVM and KSOS. SACDIN and PSOS were expected to make it to A1 as well.

[MS] So we, the seniors in the Center, were dismayed at the problems -- and the evaluator revolt -- against those expectations. It became clear that the evaluators were going to go out of their way to find fault with Multics (and by extension with Roger!) -- though we had written in either Roger's paper or in a related paper -- that we expected B2 Multics would succumb to more than one successful penetration. The evaluators' insurrection against both KSOS and Multics resulted in my direct intervention during senior staff meetings over the prolonged delays in the evaluators coming to a consensus.

[MS] There had been all kinds of "criteria lawyer" discussions/debates on Dockmaster over the strict wording in the TCSEC and how they needed to be interpreted for the

candidate systems. There were battles over whether pages not yet allocated to any segment had to be purged when they were deallocated to satisfy the memory remenance issues for newly allocated segments had to be purged prior to being allocated or whether pages that were allocated to a segment could be purged just prior to their allocation -- yeah, I know what you're thinking! So yes, interpretations took forever, and they always ended up being resolved by some management decree.

3.3.2 Multics Evaluation Team

[WOS] Deborah Downs, Aerospace Corp, Team leader
Cornelius (Neal) Haley, NSA (I think), and now an independent contractor
Maria Pozzo (now Maria King), MITRE (she moved from MITRE to Honeywell during the evaluation)
Olin Sibert, Independent consultant, Evaluator
Eric Swenson, 1st Lieutenant, Air Force Data Services Center (who likewise moved to Honeywell)
Sammy Migues, 1st Lieutenant, Air Force Data Services Center
Grant Wagner, NSA (now head of the R-division office that brought us SE/Linux)
Virgil Gligor (then a U Maryland professor and consultant to NSA). He departed around when I joined (which I think was May 1983).

[WOS] I think there were others, but my memory is failing. Might have had one or two others from MITRE or Aerospace. I was the only one not from DoD or MITRE/Aerospace. Eric and Sammy might have come a bit later; I'm not sure when they joined, but they were certainly with us once we moved to the Pentagon.

3.3.3 Matching the Model and Interpretation

[DEB] As it happens, I was at NSA while the B2 evaluation was going on. Mario Tinto, who was running the product-evaluation division at the time, asked me to take a look at the model and model interpretation documents.

[DEB] After the week-end, I met with him and said "the bad news is that there is nothing to link the Multics design documents to the model they point to [Ed. Len and my model]. the good news is I volunteer."

[DEB] So, I read the documents they had, including some design documents I hadn't seen before and drafted up a kind of cross-reference:

for function 1, it is interpreted as rule_n;

for function 2, it is interpreted as rule_k followed by rule_z;

...

to the end.

[DEB] Mario passed it on to the Multics team. After review, they returned with some comments and criticisms:

[DEB] "Now that we see what you're doing, here are some places where you're wrong -- these functions are not visible at the TCB. And here are some additional functions that *are* visible at the TCB, but that we hadn't included before."

[DEB] We went back and forth a couple of times till the designers and I were happy that we were both seeing things the same way.

3.4 Evaluation process, iterations

XXX Honeywell FSO, Cambridge and Phoenix.
NCSC seniors and evaluation team.

[WOS] I wasn't present for the beginning of the evaluation, and don't know much before then. I think it was mostly administrative activity, getting set up (with FSD) for training, getting an initial document set, etc. All the NCSC interactions were taking place on MIT-Multics at that time, but there wasn't a lot of activity.

[WOS] My involvement started when I was contracted by Honeywell FSD to teach the advanced Multics courses (F80, Internals, and F88, Dump Analysis) to an audience that included the evaluation team. After that training the team (Deborah Downs, particularly) decided that they wanted to have me available for further consultation, since they really needed a Multics technical expert, and I started consulting for NSA (through Aerospace) shortly thereafter.

[WOS] One of my early contributions was explaining how Multics development worked and where the real work was done. This really turned a corner for the evaluation, because they had previously been dealing only with FSD, which really didn't have anything to do with development.

[WOS] I wasn't really involved with the Honeywell side of the evaluation--although I was still doing some consulting work for Honeywell, it wasn't particularly evaluation-related. I recall Keith Loepere being the technical lead for evaluation work.

[WOS] About penetration testing, it was a fairly long effort: bursts of activity separated by long intervals doing other things (or working on other projects--evaluation of any one system was never a full-time activity for anyone). We didn't really have "rounds", and we didn't re-test explicitly. I don't believe the flaws ended up as part of the Honeywell-developed test suite. The five critical flaws were discovered during that effort.

[WOS] I don't believe we evaluators ever considered the successful penetrations a serious issue for Multics: they were relatively simple issues, easily remedied, and easily addressed for future work by improved coding/audit practices... which CISL/PMDC were eager to adopt, because security flaws are embarrassing. They certainly weren't fundamental design flaws.

[WOS] What *were* serious issues were things that were just missing, like:

- auditing for successful operations (fixed with new logging mechanism and lots of audit messages)
- auditing of administrative actions (partly fixed with new audit messages and better console logging)
- enforced separation of operator/administrator roles (improved, can't remember details)
- authentication for operators (fixed: operator login prompt)
- comprehensive design documentation (improved with MDDs, not completed)
- security functional testing (don't remember what happened here, but there was a huge difference between the System-M acceptance tests and what the TCSEC called for, and I think substantial work was done to address this area)
- configuration management mechanisms (fixed(?) with enhanced change comments and CM database)

[WOS] Those were all direct failures against explicit TCSEC requirements. They were pretty hard to sweep under the rug. And the news wasn't quite as welcomed by Honeywell, because the system had been pretty successful without them.

[WOS] We made a lot of compromises to get the evaluation finished, and relied a lot on promises of future work, especially for testing and documentation. That work of course never happened, since the product was canceled. That list probably isn't comprehensive--I think there were also some esoteric object security policy issues, a bunch of covert channels (Keith Loepere did a lot of work there), and various minor deficiencies both in code and documentation.

[WOS] It's interesting to hear about the "evaluator's revolt". As an evaluator, I think what we faced was a disconnect between what Roger remembered from a decade earlier and what was actually present in the product. Sure, the basic processes and segments worked fine, but there was a lot more to the system than that. This was an issue with many evaluated systems--senior staff would have, or obtain, a very rosy picture of the system, and then act surprised when that didn't match the reality. I remember some of those discussions--it was often a surprise to the senior folks when we'd describe deficiencies, and it was often a surprise to the evaluators (and developers) when the TCSEC called for stuff that just wasn't done. I mean, they gave

us the book, and if they'd wanted it to mean something different, they should have written it differently.

3.4.1 Development Process

XXX Need lots here.

Who did what.

There must have been internal task lists and progress reports.

Dev process was enhanced with extra audit.

MTB-712 and MTB-713 describe the standard process, but not the extra B2 audit.

They mention some MABs we don't have.

MTB-716 describes configuration management and enhanced history comments in the source.

Were MTBs/MABs given to the evaluation team as evidence?

3.4.2 Documentation

XXX Some Multics documentation relevant to B2 existed and had to be identified. A lot more was required by the TCSEC, see the MDD series. Internal MTBs describing plans were written and discussed according to the standard Multics process.

Presumably there were MCRs for each change, wish we had them.

3.4.3 Configuration management

XXX See MTB-716

3.4.4 Audit and Testing

XXX Separate B2 audit of all changes. Who did it. Did it find anything. Functional tests in MDD-004. Five day test run on dedicated system in PMDC.

3.4.5 Penetration Testing

XXX Done by evaluators.

What was the protocol for reporting and acting on problems.

describe -- was this all before CISL started work, were there pen tests that did not find holes

did the evaluators do only one round of pen testing? did they retest to see if holes were closed?

3.4.6 Five Critical Flaws

XXX describe -- were these all results of pen testing?

3.4.7 Covert channel analysis

XXX described in [MTB-696](#) and in [Loepere85] seems like they didn't actually address the timing channel in [Timing Channels](#).

[GMP] I recall working on covert channel analysis of the ring 1 message segment facility. I also remember Benson Margulies and I deciding to completely rewrite the message segment facility in order to properly isolate its access checking to a single module. We also added the required auditing messages at the time (see [MTB-685](#)). (Benson and I audited the original changes and determined that there were too many problems with it to fix them.)

XXX Covert channels that were not written up?

3.4.8 Technical changes made to Multics

XXX Need lots here... hoping Olin has a list

3.4.9 Anecdotes

[WOS] There are some mildly amusing stories around penetration testing. Deborah prepared a fabulous set of background material and training in the Flaw Hypothesis Methodology--this really focused our thinking [Weissman95, IBM76]. We absorbed all that, then spent a day in a windowless conference room in the (now demolished) MITRE B building. Most of the hypotheses were really generic, because at that time, only I knew very much about the system, and I expressed the opinion that while we might find some minor issues, I really didn't expect us to find any serious flaws. That night, however, I came up with the first of the five critical flaws, coded up a proof-of-concept (without crashing MIT, though there were some embarrassing SYSERR messages), and brought a demonstration the next morning. It was an eye-opener for all of us.

[WOS] We were at MITRE for a week, and then we moved to the Pentagon basement a few weeks later. Eric and Sammy had arranged for a complete set of hardcore listings to be printed, and we reviewed them page by page (in another of course windowless room) looking for time-of-check-to-time-of-use bugs and other parameter errors. Found some, too. We tested these on the AFDSC Test system, which we did crash more than once.

[WOS] During all this, I created a little application (`lister` and `compose`) to manage all the flaw information. We used that to produce very nice penetration testing reports,

with summaries and cross-references. It's possible that I still have a copy of one in my files.

[WOS] The Multics evaluation also produced the compose-based evaluation report templates that we used for a while until that was replaced by a LaTeX version some years later. Grant Wagner was the major contributor there. Jim Homan was instrumental in making LaTeX (and especially the PostScript translation) work on Multics. Neither the compose-based version nor the LaTeX version was ever really satisfactory, and it was always a black art to get the reports formatted... but it was better than the Xerox Star workstations on which a lot of the other NCSC documents were produced.

[MP] I was also involved in B2 certification. Although I can't actually recall the interfaces into the TCB I was responsible for testing I do recall that once we started to test the ring 1 mail subsystem we realized it was nowhere near B2 compliance and we rewrote it. My favorite part of B2 was covert channel analysis. I wrote a tool that would transmit data from upper level AIM classes (e.g., Top Secret) to lower classes (e.g., Unclassified) at speeds up to 110 baud. Basically, I found a way to write-down. Keith Loepere wrote a general purpose covert channel attenuator that actually slowed Multics down if one of a set of defined covert channels was detected as being exploited.

3.5 Award

XXX Who exactly decided, and how.

Letter. [\[page 1\]](#) [\[page 2\]](#)

[WOS] The Multics B2 evaluation (like all the NSA's evaluations under that regime) was divided into two phases, but there was no notion of a "provisional" certificate -- exiting the "preliminary phase" really just meant that the evaluation team had demonstrated (to an NSA internal review board) that they understood the system well enough to do the real work of the evaluation. In theory, that also meant that the product was frozen and would be evaluated as-is, but in practice (for Multics and pretty much every other evaluation) changes kept happening as more issues were discovered. Indeed, if memory serves, I believe that Multics got its B2 evaluation certificate (in July, 1985) based in part on a promise that a number of additional documents would be finished promptly thereafter.

XXX Weaknesses and issued papered over.

3.6 Subsequent Maintenance

XXX Were subsequent changes to Multics re-evaluated, and how? Was Multics part of RAMP?

3.7 Other

XXX MLS TCP/IP in ring 1. (Ring 1 not subject to AIM rules)
MRDS and forum ran in ring 2
ring 3 for site secure software. e.g. Roger's pun of the day at MIT.

[GMP] The Multics Mail System ran in ring 2. ([MTB-613](#) documents the mail_system_ API.)

XXX IP security option sent AIM bits over the network.. was this used for anything at AFDSC or elsewhere?

[WOS] AIM did start with 18 orthogonal categories; the folks who ran the NSA's Dockmaster Multics later changed it to 9 orthogonal categories plus one of 510 extended categories.

[GMP] When I created the Inter-Multics File Transfer (IMFT) system, I had it handle AIM by mapping categories and levels between systems based on the names assigned to the categories and levels at the two sites. (I don't know if this was ever actually used by a customer. I'm pretty sure the CISL/System-M connection didn't care.) I did give a HLSUA talk (in Detroit) on IMFT which might have some details on its handling of AIM.

4. Results

At the beginning of the 1970s, not many people were interested in computer security. There was no clear understanding of what was wanted. There was no commonly understood methodology for making secure systems.

By 1986, the OS security community had expanded. We had OS security conferences, polyinstantiation, and DOCKMASTER. One system, SCOMP, had been rated A1, and Multics had been awarded a B2 rating, and had been **canceled by Honeywell Bull**.

Design was begun at Bull for a follow-on product named Opus, targeted for a new machine. The target evaluation level for Opus was B3. This product was also canceled by Bull in 1988.

By 1990, no additional operating systems had been rated A1, B3, or B2.

4.1 Good

The B2 effort changed our understanding of computer security from wishful thinking and "we tried really hard" to a more rigorous analysis. It proposed (and exposed problems with) evaluation methodology. The early claims Multics made for security were shown to be insufficient, and were superseded by a more thorough model of security and a more concrete understanding of the security that Multics provided. The B2 process was also a source of security corrections to Multics. It may have kept Multics alive for a few more years. Some Multics systems may have been sold due to the (promise of future) rating.

4.2 Multics systems that used B2 features

XXX What and how.

- AFDSC (2 of the 6 systems were multilevel)
- DOCKMASTER, no levels but used categories
- Site N (Flagship) probably
- DND-H, DND-DDS
- MDA-TA (SEDACS)
- NWGS?
- RAE, first UK site to run MLS
- SEP?
- SNECMA?
- ONERA?
- CNO source control
- Oakland University used AIM to keep the students out of the source code.
- others, e.g. GM/Ford?

4.3 Problems

XXX Pressure on NSA and evaluators to award B2 rating, to prove B2 is real.

Pressure to be objective and fair.

Never done this before.

Many statements in Orange Book required interpretation.

Things not done or skated over.

Saltzer's list of areas of concern.

Things outside scope, e.g. network security, multi-system networks, heterogeneous networks

4.4 In what sense Multics was "secure"

Multics was a "securable" mainframe system. It was designed to be run in a closed machine room by trained and trusted operators, and supported by a trusted manufacturer. The threat environment, at the time that Multics was designed, had no public examples of hostile penetration. The system's design also assumed that some security decisions were not made automatically, but were left to the skill of a Site Security Officer. For instance, users of Compartmented Mode Workstations and similar systems in the 90s had to deal with downgrading problems: we left all those to the SSO.

4.5 Relevance to current and future security needs

The high-assurance, military security, mainframe oriented model of the Orange Book is oriented toward the problems of the 1970s. We assumed that Multics systems ran on trusted hardware that was maintained and administered by trained and trusted people. People didn't buy graphics cards on Canal Street and plug them into the mainframe. In many ways, our security design pushed a lot of hard-to-solve issues off onto Operations and Field Engineering. Multics hardware originated from a trusted manufacturer and was shipped on trusted trucks, installed by trusted FEs, and kept in a locked room.

The Multics B2 experience influenced the plans for Digital's A1-targeted VAX/SVS development. [Lipner12] That product was cancelled because not enough customers wanted to buy it, and because the expected development and evaluation cost of adding network support and graphical user interfaces would have been too high.

[SBL] We also know that the models for OS security really aren't great as yet. Bell and LaPadula produces an unusable system - see the MME experience with message handling. Type enforcement is probably better but dreadfully hard to administer as I recall. Nice for static systems.

Different security models are appropriate for different business needs. The government MLS classification model is one market; Multics business customers struggled to apply it to their needs. The Biba data integrity model was implemented in VAX/SVS but never used, according to Steve Lipner. The Clark-Wilson security model was oriented toward information integrity for commercial needs; no high-assurance systems have implemented this model.

Multics influence continues to affect secure development. Recent papers by Multicians on multi-level security and smartcard operating systems continue to build on the lessons we learned. [Cheng07, Karger10] Multics influence is cited in the plans for secure development efforts such as CRASH.

The Orange Book assurance process was too costly per operating system version. *Computers At Risk* [NRC91] talks about how the Orange Book bundles functional levels and level of assurance. It describes how the Orange Book process evolved into ITSEC, which provided more flexibility in defining what functions to assure, and also mentions issues with the incentives for system evaluators.

4.6 Conclusion

What we learned from Roger Schell:

1. Security through obscurity doesn't work. Attackers can discover obscure facts. Any security weakness, no matter how obscure, can be exploited. In particular, keeping the source code secret doesn't work. Attackers get hold of source code, or do without.
2. Penetrate and patch doesn't work. You're never sure that an attacker has not found a new vulnerability that your penetrators didn't.
3. Counting on attackers to be lazy doesn't work. They will expend very large amounts of time and effort looking for weaknesses and inventing exploits, much more than simple economic models justify.
4. Security has to be formally defined relative to a policy model. The model has to be consistent: it can't lead to contradictions.
5. To demonstrate that a system is secure requires a logical argument showing the interpretation of the policy model with respect to the system design.
6. The part of a system that makes access control decisions is the *reference monitor*.
7. Showing that a reference monitor correctly implements a security model requires a logical demonstration. To do this, the monitor must be small enough to analyze completely (a *security kernel*). Unless the monitor is very small, this requires automated analysis.
8. Proofs of security must still be verified by thorough testing.
9. System security has to be supported through the whole system life cycle, including distribution and maintenance.

After the B2 exercises, do we now know how to build a "secure" OS? We know measures that help; we know things to avoid doing, including security-by-wishful-thinking. The few worked examples come with lots of qualifications and excuses.

Was it worth it? The marketplace seems to have answered this question: not at current prices.

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